

# Mechanisation of the AKS Algorithm

Hing-Lun Chan

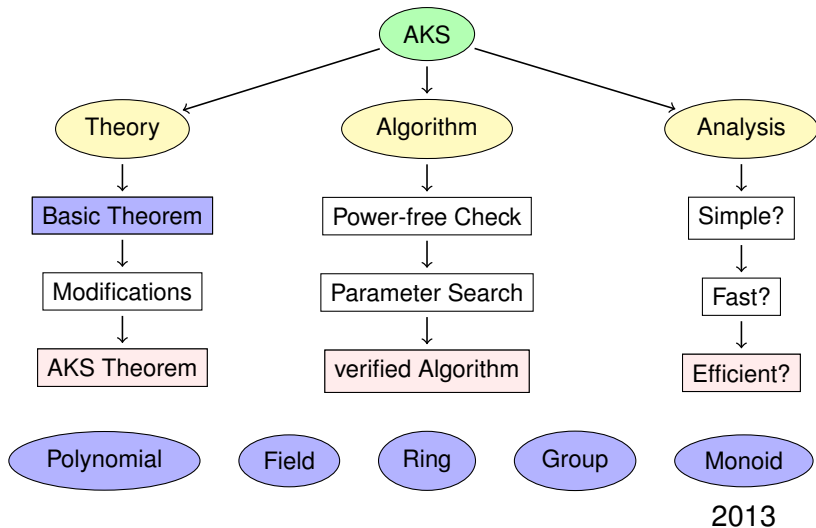
College of Engineering and Computer Science  
Australian National University

PhD Review 2016

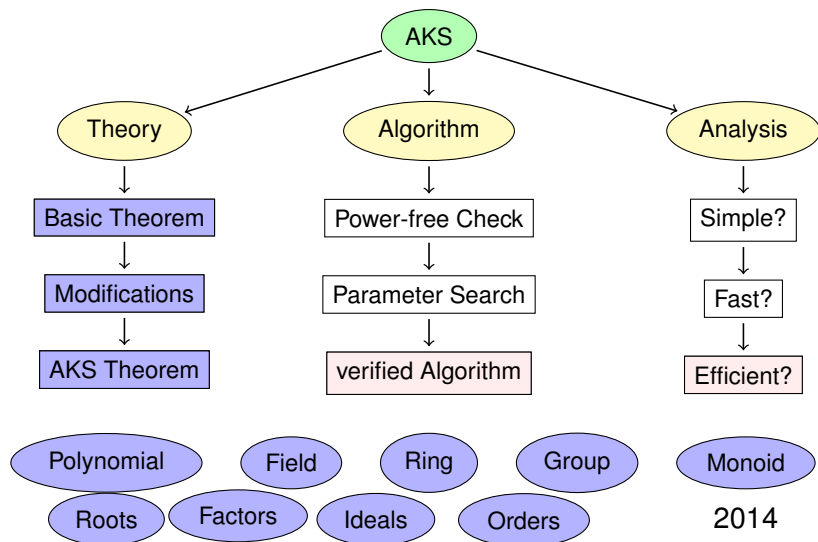
# Outline

- 1 Introduction
  - Road Map
- 2 Work Progress
  - View in 2014
  - View in 2015
  - View in 2016
- 3 Big Picture
  - Jigsaw Puzzle
- 4 Look Ahead
  - Plans

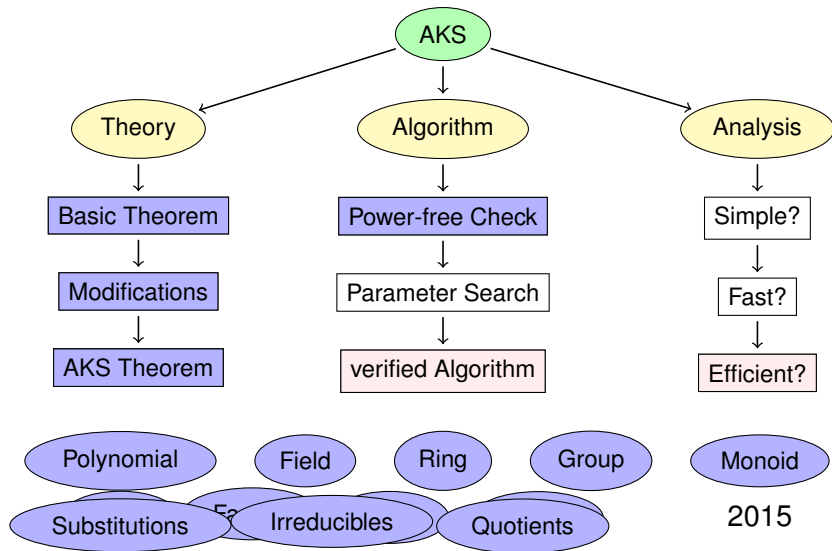
# Mechanisation of AKS Algorithm – Road Map



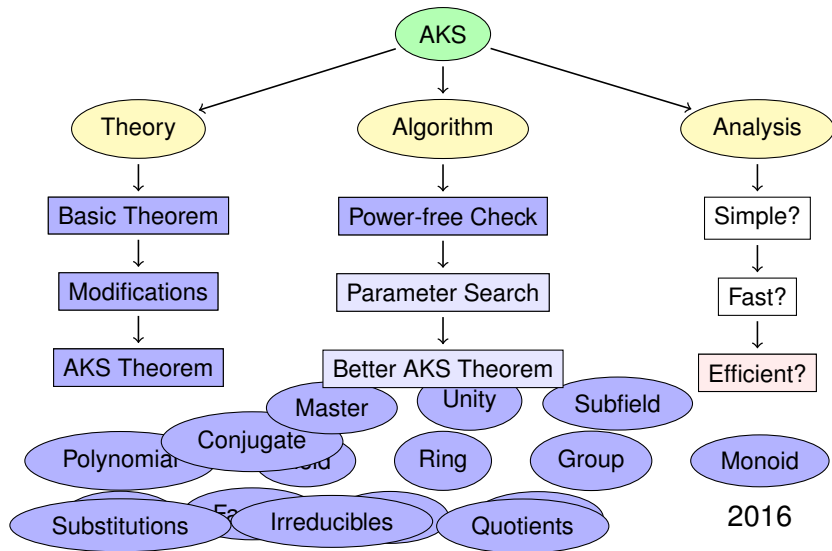
# Mechanisation of AKS Algorithm – Road Map



# Mechanisation of AKS Algorithm – Road Map



# Mechanisation of AKS Algorithm – Road Map



# Back in 2014

## Completed:

- Basic libraries for Group, Ring, Field.
- Basic libraries for Polynomials.
- Show that Polynomials modulo an irreducible form a finite field.

## To Do:

- Understand the role of [parameters](#) for the AKS algorithm.
- Formulate the notion of [introspective](#) relation in AKS proof.
- Formulate the sets and maps critical for the AKS proof.
- How the Pigeonhole Principle will play in AKS proof?

# Back in 2015

## Completed:

- Put AKS in script, with one **parameter**  $k$ , derived from input  $n$ .
- Correct definition of **introspective** relation after some false starts.
- AKS parameter  $k$  imposes injective maps between sets.
- Key: if AKS is false, an injective map between finite sets would violate the Pigeonhole Principle.
- Part 1: a prime  $k \Rightarrow$  AKS Main Theorem, and such  $k$  exists.

## To Do:

- Why 2nd version of AKS proof does not require  $k$  to be prime?
- **Remove** the prime requirement on  $k$  for the formalized AKS proof.
- Show that parameter  $k$  can be found within steps of  $O(\log n)$ .



# Now in 2016

## Revised:

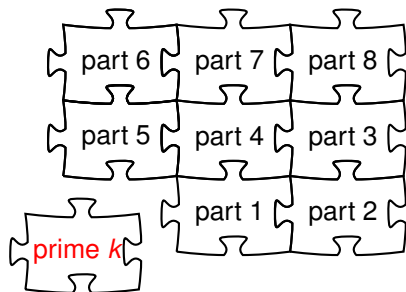
- Need: to gather information about Cyclotomic factors.
- Need: to establish the existence of Finite Fields.
- Key: obtain a count of monic irreducibles for a given degree.
- Todo: Reformalute AKS proof with a simple parameter  $k$ .

## Achieved:

- Search for simple  $k$  is  $O(\log n)$  if an LCM bound is true.
- A short joint paper to ITP2016 for a cute proof of this result:  
 $2^n \leq LCM \{1; 2; 3; \dots; (n + 1)\}$ .
- An implicit formula for the monic irreducibles count.
- A finite field exists with cardinality  $p^n$ , for prime  $p$  and  $0 < n$ .

# AKS Jigsaw Puzzle

Parameter  $k$  is derived from input  $n$

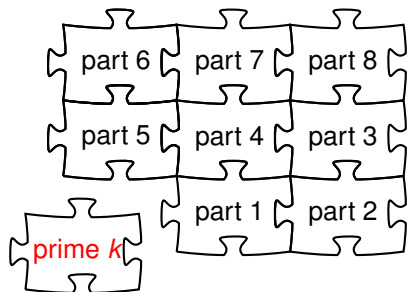


Input: a **power-free number**  $n$

- 1: Coprime checks for  $\{1; \dots; k\}$
- 2: From  $k$  derive another limit  $\ell$
- 3: Polynomial checks for  $\{1; \dots; \ell\}$   
in  $(\text{mod } n, X^k - 1)$
- 4: Choose  $p$ , a prime divisor of  $n$
- 5: Shift to  $(\text{mod } p, X^k - 1)$
- 6: Construct introspective sets
- 7: Craft finite sets and 1-to-1 maps
- 8: Conclude:  $n$  must be a power of  $p$   
or Pigeonhole principle is violated!

# AKS Jigsaw Puzzle

Parameter  $k$  is derived from input  $n$



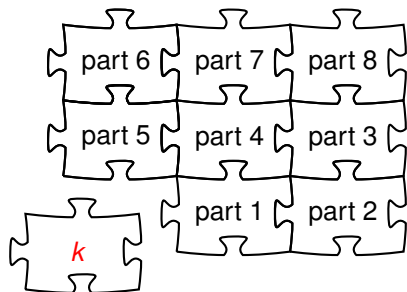
Input: a power-free number  $n$

- 1: Coprime checks for  $\{1; \dots; k\}$
- 2: From  $k$  derive another limit  $\ell$
- 3: Polynomial checks for  $\{1; \dots; \ell\}$   
in  $(\text{mod } n, X^k - 1)$
- 4: Choose  $p$ , a prime divisor of  $n$
- 5: Shift to  $(\text{mod } p, X^k - 1)$
- 6: Construct introspective sets
- 7: Craft finite sets and 1-to-1 maps
- 8: Conclude:  $n$  must be a power of  $p$   
or Pigeonhole principle is violated!

A prime  $k$  makes a key proof step easy.

# AKS Jigsaw Puzzle – improved

Parameter  $k$  is derived from input  $n$

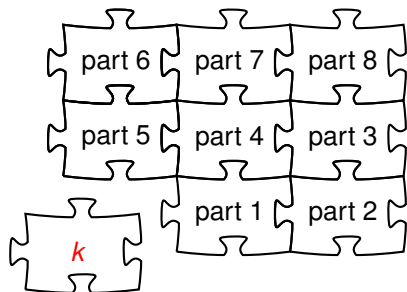


Input: a **power-free number**  $n$

- 1: Coprime checks for  $\{1; \dots; k\}$
- 2: From  $k$  derive another limit  $\ell$
- 3: Polynomial checks for  $\{1; \dots; \ell\}$   
in  $(\text{mod } n, X^k - 1)$
- 4: Choose  $p$ , a prime divisor of  $n$
- 5: Shift to  $(\text{mod } p, X^k - 1)$
- 6: Construct introspective sets
- 7: Craft finite sets and 1-to-1 maps
- 8: Conclude:  $n$  must be a power of  $p$   
or Pigeonhole principle is violated!

# AKS Jigsaw Puzzle – improved

Parameter  $k$  is derived from input  $n$



Input: a **power-free number**  $n$

- 1: Coprime checks for  $\{1; \dots; k\}$
- 2: From  $k$  derive another limit  $\ell$
- 3: Polynomial checks for  $\{1; \dots; \ell\}$   
in  $(\text{mod } n, X^k - 1)$
- 4: Choose  $p$ , a prime divisor of  $n$
- 5: Shift to  $(\text{mod } p, X^k - 1)$
- 6: Construct introspective sets
- 7: Craft finite sets and 1-to-1 maps
- 8: Conclude:  $n$  must be a power of  $p$   
or Pigeonhole principle is violated!

Just  $k$  will simplify its search bound.

# Possible Timeline

Thesis plan:

**April, 2015:** AKS Main Theorem (with suitable prime  $k$ )

**June, 2016:** AKS Main Theorem (with suitable  $k$ )

**June, 2017:** Complexity/Efficiency of AKS algorithm

**December, 2017:** Thesis written (hopefully!)

# Possible Timeline

Thesis plan:

**April, 2015:** AKS Main Theorem (with suitable prime  $k$ )

**June, 2016:** AKS Main Theorem (with suitable  $k$ )

**June, 2017:** Complexity/Efficiency of AKS algorithm

**December, 2017:** Thesis written (hopefully!)

Publications:

**CPP2012:** A String of Pearls: Proofs of Fermat's Little Theorem

**JFR2013:** Extended version for Journal of Formalized Reasoning

**ITP2015:** Mechanisation of AKS Algorithm Part 1: Main Theorem

**ITP2016:** (submitted) Bounding LCMs with Triangles

**??** (planned) Mechanisation of AKS Algorithm Part 2.