

## 1 Expressions

$\neg p \vee p$  vs.  $\neg p \vee p$

$\neg(p \wedge \neg p)$  vs.  $\neg(p \wedge \neg p)$

$(p \rightarrow (q \wedge r)) \rightarrow (q \wedge r \wedge p)$

$(q, r) \supset (q \supset r)$

## 2 Evaluation

$$\frac{\begin{array}{c} p \rightarrow \perp \\ q \rightarrow \top \end{array}}{\begin{array}{c} p \rightarrow 0 \\ q \rightarrow 1 \end{array}} \quad \frac{}{\begin{array}{c} p \rightarrow F \\ q \rightarrow T \end{array}}$$

$$\frac{\begin{array}{c} p \rightarrow \perp \\ q \rightarrow \top \end{array}}{\begin{array}{c} p \rightarrow \perp \\ q \rightarrow \top \end{array}} \quad \frac{}{r \rightarrow T}$$

$$\frac{\begin{array}{c} p \rightarrow \perp \\ q \rightarrow \top \\ r \rightarrow \perp \\ s \rightarrow \top \end{array}}{\quad}$$

$$\frac{\begin{array}{c} p \rightarrow \perp \\ q \rightarrow \perp \end{array}}{\begin{array}{c} p \vee q \rightarrow \perp \\ p \wedge q \rightarrow \perp \\ p \rightarrow q \rightarrow \top \\ p \rightarrow (\neg q) \rightarrow \top \\ (\neg p) \rightarrow q \rightarrow \perp \\ p \wedge \neg q \rightarrow \perp \\ \neg p \wedge q \rightarrow \perp \end{array}}$$

$$\frac{\begin{array}{c} p \rightarrow F \\ q \rightarrow T \end{array}}{\begin{array}{c} p \vee q \rightarrow T \\ p \wedge q \rightarrow F \\ p \rightarrow q \rightarrow T \\ p \rightarrow (\neg q) \rightarrow T \\ (\neg p) \rightarrow q \rightarrow T \\ p \wedge \neg q \rightarrow F \\ \neg p \wedge q \rightarrow T \end{array}}$$

$$\frac{\begin{array}{c} p \rightarrow \top \\ q \rightarrow \perp \end{array}}{\begin{array}{c} p \vee q \rightarrow \top \\ p \wedge q \rightarrow \perp \\ p \rightarrow q \rightarrow \perp \\ p \rightarrow (\neg q) \rightarrow \top \\ (\neg p) \rightarrow q \rightarrow \top \\ p \wedge \neg q \rightarrow \top \\ \neg p \wedge q \rightarrow \perp \end{array}}$$

$$\frac{\begin{array}{c} p \rightarrow 1 \\ q \rightarrow 1 \end{array}}{\begin{array}{c} p \vee q \rightarrow 1 \\ p \wedge q \rightarrow 1 \\ p \rightarrow q \rightarrow 1 \\ p \rightarrow (\neg q) \rightarrow 0 \\ (\neg p) \rightarrow q \rightarrow 1 \\ p \wedge \neg q \rightarrow 0 \\ \neg p \wedge q \rightarrow 0 \end{array}}$$

a

b

c

d

### 3 Truth tables

$\wedge$	$p$	$q$
0	0	0
0	0	1
1	1	0
0	1	1

$p$	$\wedge$	$\neg$	$q$
0	0	1	0
0	0	0	1
1	1	1	0
1	0	0	1

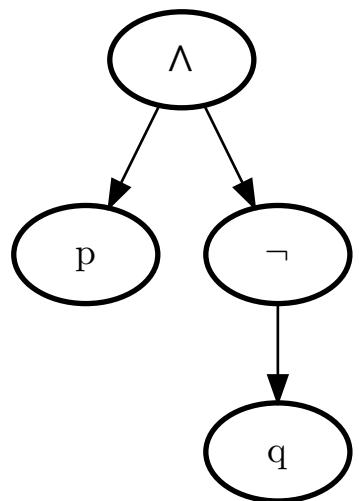
$p$	$\wedge$	$\neg$	$q$
0	0	1	0
0	0	0	1
1	1	1	0
1	0	0	1

$((p \rightarrow q) \rightarrow (r \rightarrow (p \rightarrow q) \rightarrow (r \rightarrow q)))$	$\rightarrow$	$(r \rightarrow (p \rightarrow q) \rightarrow (r \rightarrow q))$	$\rightarrow$	$(p \rightarrow q) \rightarrow (r \rightarrow q)$	$\rightarrow$	$r \rightarrow q$	$\rightarrow$	$p \rightarrow q$	$\rightarrow$	$p \rightarrow q$	$\rightarrow$	$r \rightarrow q$
0	1	0	1	0	1	0	1	0	1	0	1	0
0	1	0	0	1	0	0	1	0	0	1	0	0
0	1	1	1	0	1	0	1	1	1	0	1	1
0	1	1	1	1	1	0	1	1	1	1	1	1
1	0	0	1	0	1	1	0	0	1	0	1	0
1	0	0	1	1	1	1	0	0	1	1	0	0
1	1	1	1	0	1	1	1	1	1	0	1	1
1	1	1	1	1	1	1	1	1	1	1	1	1

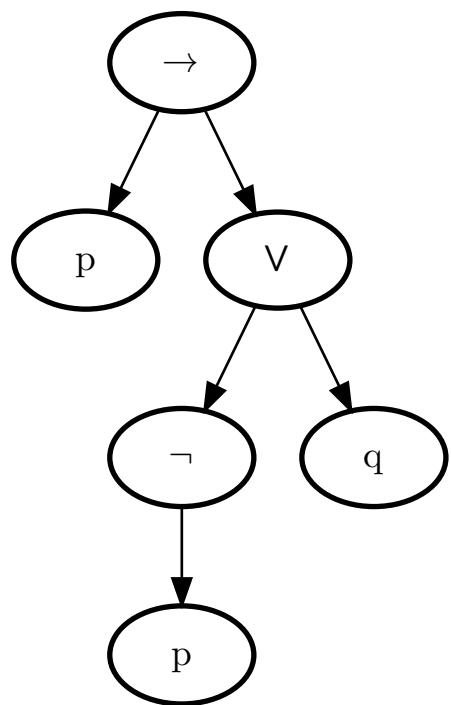
$p$	$\wedge$	$(\neg q \rightarrow r)$
0	0	1
0	0	0
0	0	1
0	0	0
0	0	1

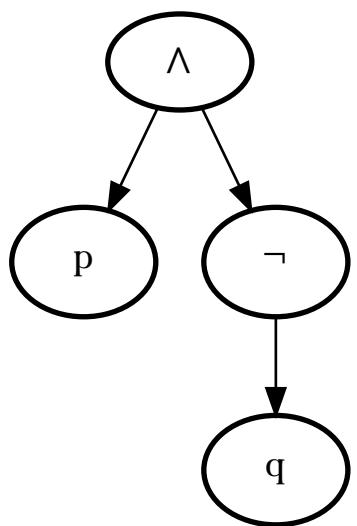
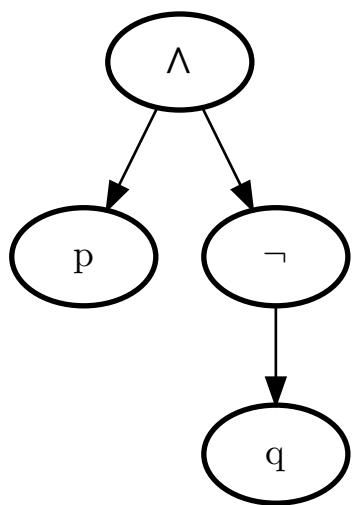
1	0	1	0	0	0
1	1	1	0	1	1
1	1	0	1	1	0
1	1	0	1	1	1

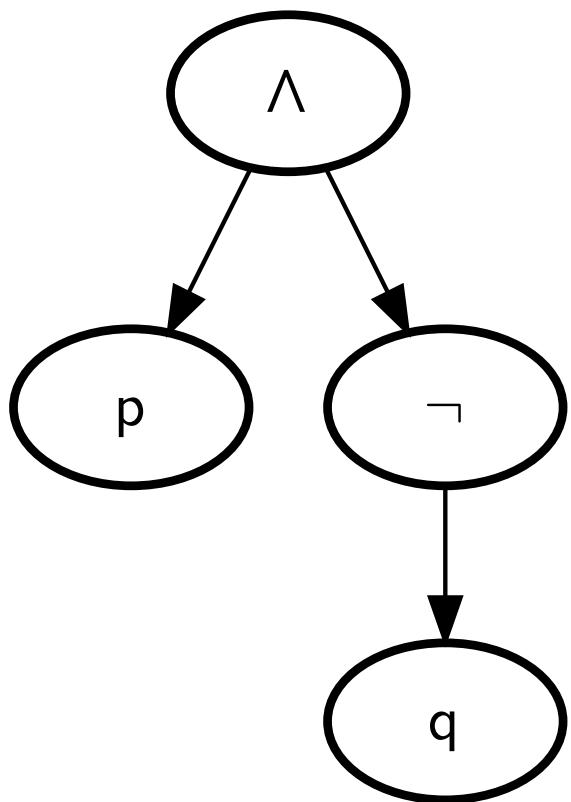
## 4 Syntax trees

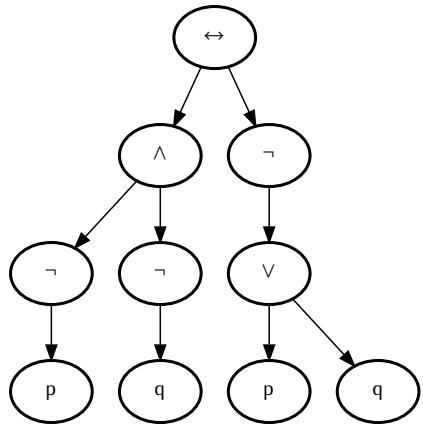


$\rightarrow$	$p$	$q$
1	0	0
1	0	1
0	1	0
1	1	1

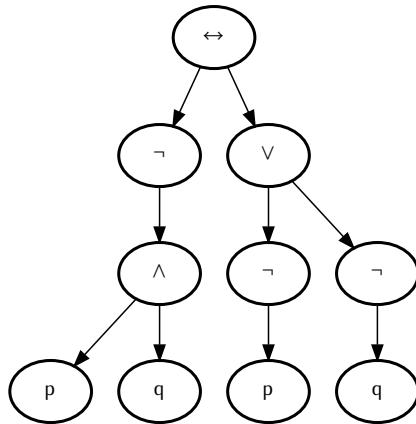








DeMorgan's law 1



DeMorgan's law 2

$\leftrightarrow$	$p$	$q$
1	0	0
1	0	1
1	1	0
1	1	1

Truth  
table 1

$\leftrightarrow$	$p$	$q$
1	0	0
1	0	1
1	1	0
1	1	1

Truth  
table 2

## 5 Venn diagrams

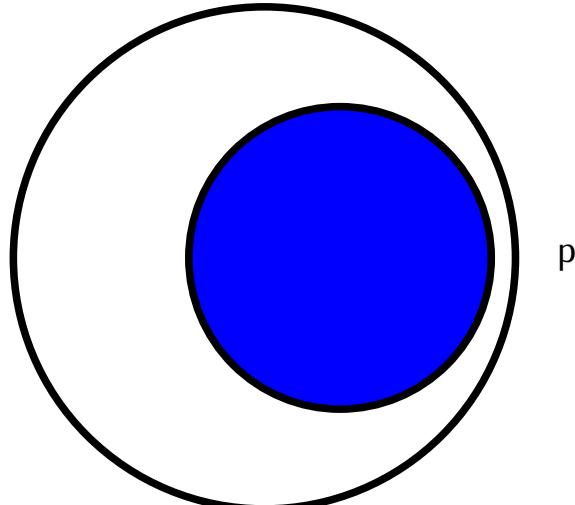


Figure 1  $\neg p$

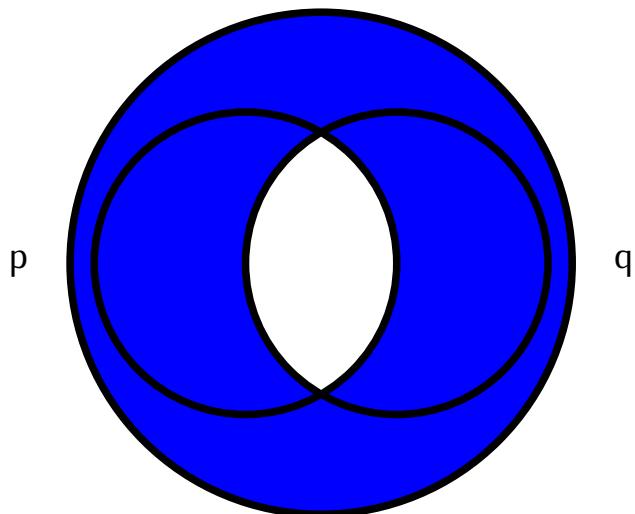


Figure 2  $p \wedge q$

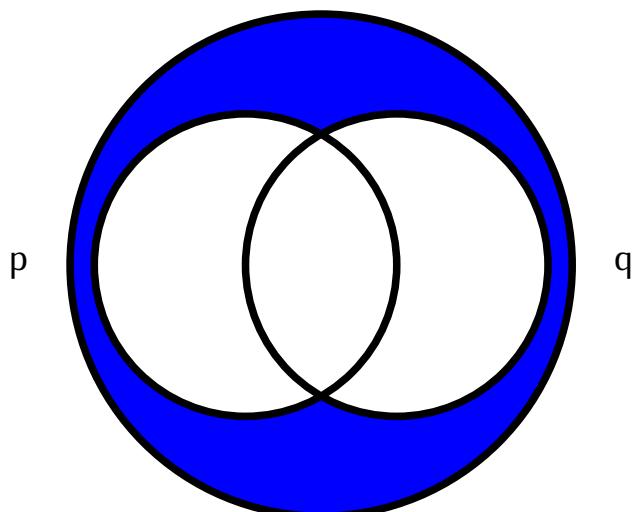


Figure 3  $p \vee q$

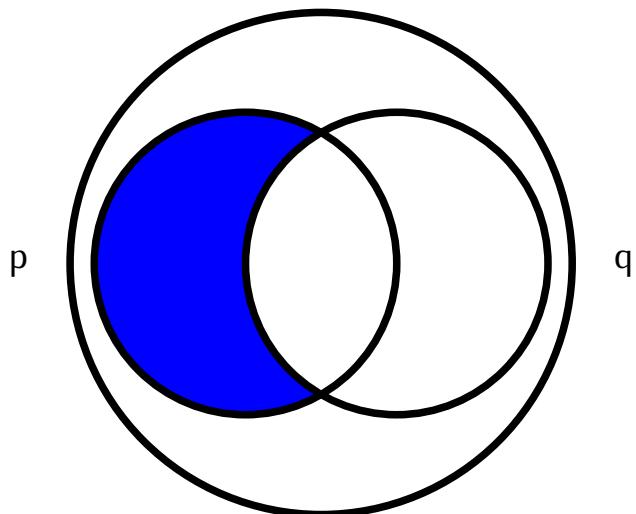


Figure 4  $p \supset q$

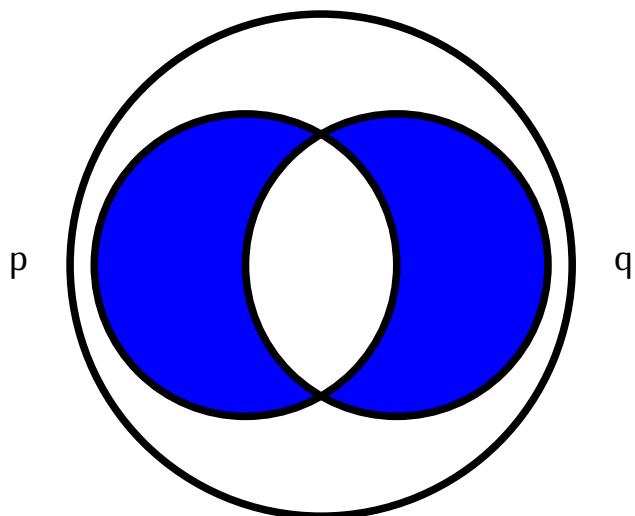


Figure 5  $p \leftrightarrow q$

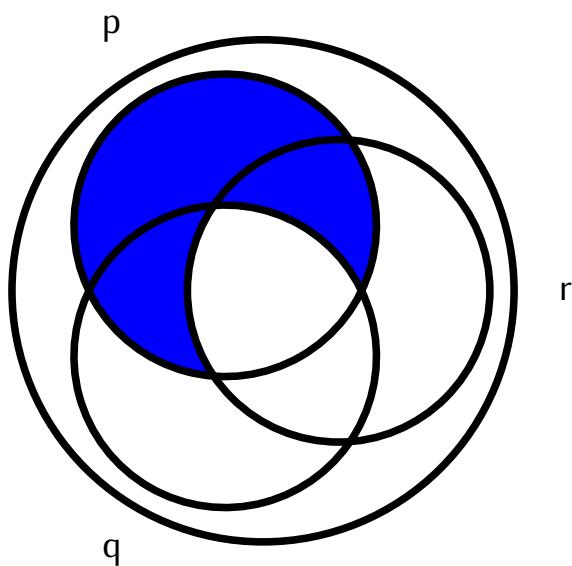


Figure 6  $p \rightarrow (q \wedge r)$

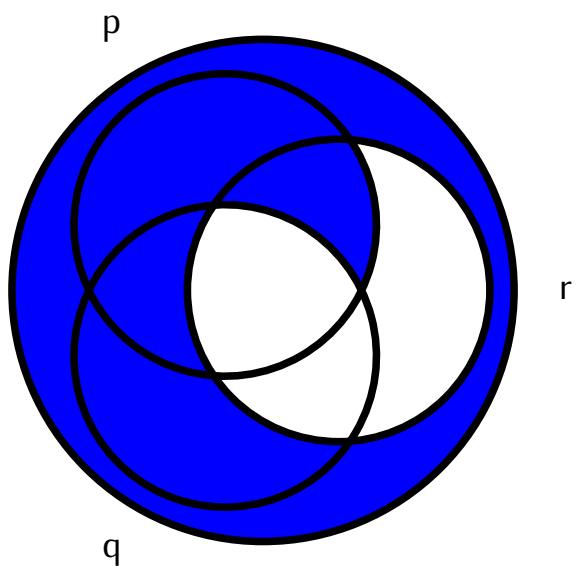


Figure 7  $(p \rightarrow q) \wedge r$